

FORM PTO-1390

U S DEPARTMENT OF COMMERCE PATENT AND TRADEMARK OFFICE

ATTORNEY'S DOCKET NUMBER

616-034

TRANSMITTAL LETTER TO THE UNITED STATES
DESIGNATED/ELECTED OFFICE (DO/EO/US)
CONCERNING A FILING UNDER 35 U.S.C. 371

U.S. APPLICATION NO. (if known, see 37 CFR 1.5)

09/743713

INTERNATIONAL APPLICATION NO.

PCT/GB99/02288

INTERNATIONAL FILING DATE

15 July 1999 (15.07.99)

PRIORITY DATE CLAIMED

16 July 1998 (06.07.98)

TITLE OF INVENTION

CONTENT ADDRESSED MEMORIES

APPLICANT(S) FOR DO/EO/US

Douglas Philip TURVEY

Applicant herewith submits to the United States Designated/Elected Office (DO/EO/US) the following items and other information:

1. ☒ This is a **FIRST** submission of items concerning a filing under 35 U.S.C. 371.
2. ☐ This is a **SECOND** or **SUBSEQUENT** submission of items concerning a filing under 35 U.S.C. 371.
3. ☐ This express request to begin national examination procedures (35 U.S.C. 371(f)) at any time rather than delay examination until the expiration of the applicable time limit set in 35 U.S.C. 371(b) and PCT Articles 22 and 39(1).
4. ☐ A proper Demand for International Preliminary Examination was made by the 19th month from the earliest claimed priority date.
5. ☐ A copy of the International Application as filed (35 U.S.C. 371(c)(2))
 - a. ☐ is transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ has been transmitted by the International Bureau.
 - c. ☐ is not required, as the application was filed in the United States Receiving Office (RO/US)
6. ☐ A translation of the International Application into English (35 U.S.C. 371(c)(2)).
7. ☐ Amendments to the claims of the International Application under PCT Article 19 (35 U.S.C. 371(c)(3))
 - a. ☐ are transmitted herewith (required only if not transmitted by the International Bureau).
 - b. ☐ have been transmitted by the International Bureau.
 - c. ☐ have not been made; however, the time limit for making such amendment has NOT expired.
 - d. ☐ have not been made and will not be made.
8. ☐ A translation of the amendments to the claims under PCT Article 19 (35 U.S.C. 371(c)(3)).
9. ☐ An oath or declaration of the inventor(s) (35 U.S.C. 371(c)(4)).
10. ☐ A translation of the annexes to the International Preliminary Examination Report under PCT Article 36 (35 U.S.C. 371(c)(5)).

Items 11. To 16. below concern other document(s) or information included:

11. ☐ An Information Disclosure Statement under 37 CFR 1.97 and 1.98.
12. ☐ An assignment document for recording. A separate cover sheet in compliance with 37 CFR 3.28 and 3.31 is included.
13. ☒ A **FIRST** preliminary amendment.
14. ☐ A **SECOND** or **SUBSEQUENT** preliminary amendment.
15. ☐ A substitute specification.
16. ☐ A change of power of attorney and/or address letter.
17. ☐ Other items or information.

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- 17.
- ☒
- The following fees are submitted:

Basic National Fee (37 CFR 1.492(a)(1)-(5)):

Neither international preliminary examination fee (37 CFR 1.482) nor international search fee (37 CFR 1.445(a)(2)) paid to USPTO and International Search Report not prepared by the EPO or JPO \$ 1,000.00

International preliminary examination fee (37 CFR 1.482) not paid to USPTO but International Search Report prepared by the EPO and JPO \$ 860.00

International preliminary examination fee (37 CFR 1.482) not paid to USPTO but international search fee (37 CFR 1.445(a)(2)) paid to USPTO \$ 710.00

International preliminary examination fee paid to USPTO (37 CFR 1.482) but all claims did not satisfy provisions of PCT Article 33(1)-(4) \$ 690.00

International preliminary examination fee paid to USPTO (37 CFR 1.482) And all claims satisfied provisions of PCT Article 33(2)-(4) \$ 100.00

ENTER APPROPRIATE BASIC FEE AMOUNT =

CALCULATIONS PTO USE ONLY

\$ 860.00

Surcharge of \$130.00 for furnishing the oath or declaration later than ☐ 20 ☒ 30 months from the earliest claimed priority date (37 CFR 1.492(e)).

\$ 130.00

CLAIMS	NUMBER FILED	NUMBER EXTRA	RATE		
Total Claims	11 - 20 =		X \$18.00	\$	0.00
Independent Claims	2 - 3 =		X \$80.00	\$	0.00
Multiple dependent claim(s) (if applicable)			+ \$260.00	\$	0.00

TOTAL OF ABOVE CALCULATIONS =

\$ 990.00

Reduction by 1/2 for filing by small entity, if applicable. Verified Small Entity Statement must also be filed. (Note 37 CFR 1.9, 1.27, 1.28).

\$ 495.00

SUBTOTAL =

\$ 495.00

Processing fee of \$130.00 for furnishing the English translation later than the ☐ 20 ☐ 30 months from the earliest claimed priority date (37 CFR 1.492(f)).

\$ 0.00

TOTAL NATIONAL FEE =

\$ 495.00

Fee for recording the enclosed assignment (37 CFR 1.21(h)). The assignment must be accompanied by an appropriate cover sheet (37 CFR 3.28, 3.31). \$40.00 per property

\$ 0.00

TOTAL FEES ENCLOSED =

\$ 495.00

Amount to be:
Refunded

\$

Charged

\$

- a. ☒ Please charge my Credit Card in the amount of \$ 495.00. A Credit Card Payment Form (PTO-2038) is enclosed.
- b. ☐ Please charge my Deposit Account No. XXX in the amount of \$ XXX to cover the above fees. A duplicate copy of this sheet is enclosed.
- c. ☒ The Commissioner is hereby authorized to charge any additional fees which may be required, or credit any overpayment to Deposit Account No. 07-1337. A duplicate copy of this sheet is enclosed.

NOTE: Where an appropriate time limit under 37 CFR 1.494 or 1.495 has not been met, a petition to revive (37 CFR 1.137(a) or (b)) must be filed and granted to restore the application to pending status.

SEND ALL CORRESPONDENCE TO:

Allan M. Lowe
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SIGNATURE

Allan M. Lowe

NAME

19,641

REGISTRATION NUMBER

09/743713

JC07 Rec'd POW/PIS 16 JAN 200

Docket No.: 616-034

PATENT

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

In re Application of :
Douglas Philip TURVEY :
Serial No. Not Yet Assigned :
Filed: January 16, 2001 :

For: CONTENT ADDRESSED MEMORIES

PRELIMINARY AMENDMENT

ASSISTANT COMMISSIONER FOR PATENTS
Washington, D. C. 20231

Sir:

Prior to examination on the merits, please amend the referenced application as follows:

IN THE CLAIMS:

In claim 3, line 1, please change "either previous claim" to "claim 1"

In claim 6, line 1, please change "any previous claim" to "claim 1"

In claim 7, line 1, please change "any previous claim" to "claim 1"

In claim 8, line 1-2, please change "any previous claim" to "claim 1"

In claim 9, line 1-2, please change "any previous claim" to "claim 1"

In claim 10, line 1-2, please change "any previous claim" to "claim 1"

In claim 11, line 1-2, please change "any previous claim" to "claim 1"

Please delete claim 12

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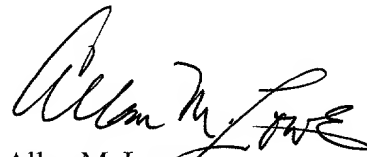
Serial No. Not Yet Assigned

Remarks

The above-referenced application is amended to delete the multiple dependencies of claims 3, and 6-11 to avoid the multiple dependent claim filing fee.

Respectfully submitted.

LOWE HAUPTMAN GILMAN & BERNER, LLP

A handwritten signature in black ink, appearing to read "Allan M. Lowe", is written over the printed name.

Allan M. Lowe
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16 JAN 2001

Content Addressed Memories

The present invention relates to memories of the type known as CAMs (Content Addressed Memories) or Associative Memories.

A conventional computer memory consists of a large number of memory locations which have sequential addresses. To access a location in such a memory, an address is supplied to the memory; the corresponding location is thereby selected, and its contents can then be accessed (ie read or written).

This works admirably when the address of the desired location (which typically stores a byte or a word) is known. However, there is a variety of situations in which data is organized in data blocks and what is wanted is a data block having particular data in a part of the block. Conventionally, the only way of finding the desired data block is to search through the data blocks one by one. A simple sequential search is the simplest procedure, but is liable to be extremely time-consuming. Various forms of directory structures or indexing can sometimes be used, but these tend to be complicated and inflexible.

To overcome these problems, a type of memory known as content addressable memory (CAM) has been proposed. This type of memory is also known as associative, because it automatically associates the desired data with the blocks containing that data.

CAMs have never achieved substantial commercial success, primarily because of the complexity of the circuitry required, compared with conventional memories. The situation is made worse because of the relatively limited range of applications for CAMs, so that CAM memories would be a relatively low-volume and high cost product even apart from the extra complexity.

The general object of the present invention is to provide an improved CAM architecture which alleviates or overcomes these problems.

According to the present invention there is provided a content addressable memory comprising a CAM control logic unit and a plurality of cells connected in a chain, each cell comprising:

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a memory block coupled to a common address bus;

a comparator coupled to a common data bus and to the data interface of the memory block;

switching means coupling the data interface of the memory block with the data bus, and;

a logic block including a Match flip-flop:

the memory being operable:

in a Search phase to serially match a sequence of words on the common data bus with the contents of a sequence of addresses in the memory blocks of the cells; and

in an Access phase, to render the cells matched in the Search phase serially available for access via the common address and data buses.

A CAM embodying the invention will now be described, by way of example, with reference to the drawings, in which:

Fig. 1 is a general block diagram of the system;

Fig. 2 is a more detailed block diagram of a cell of the system; and

Fig. 3 is a detailed block diagram of the logic block of a cell of the system.

Referring to Fig. 1, the CAM consists primarily of a chain of identical cells 10, with a CAM control unit 11, which is coupled to the top of the chain of logic blocks 10.

Each cell (Fig. 2) contains a respective memory block 12, which may use say DRAM (with conventional refresh arrangements operating on each block) or SRAM. The size of each memory block may be say 64 kbytes, though obviously the number of locations may vary and the word length may also be varied (eg to 2 or 4 bytes). The number of blocks will normally be large, typically of the order of 105.

The cells 10 are coupled to a 1-byte data bus DATA, a 16-bit address bus ADD, and a control bus CONT. The address and data bus widths match the number of locations in the memory block in each cell and the size of each location. If the CAM forms part of a computer system, these buses may be general system-wide buses, or parts of such system buses. A system address bus may be consi-

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derably wider than 16 bits; the bottom 16 bits will then be used as the address bus ADD for the CAM. Similarly, the control bus for the CAM may form part of a multi-bit system control bus.

Each cell 10 also contains a memory block 12, a logic block 13, a comparator 14, and a bidirectional switch 15. The chaining of the cells 10 is through their logic blocks 13, as shown.

The CAM will normally be implemented on an integrated circuit chip. It will be realized that such a chip can easily be designed so that several such chips can be chained to increase the size of the CAM. The control unit 11 is preferably a separate unit. However, it can be included on the CAM chip if desired; in that case, several chips can be chained by designing the controller so that only that of the first chip in the chain is enabled, with those of the following chips in the chain being disabled.

The functions of the logic blocks 13 can best be understood by considering the manner in which the CAM operates. The CAM operates in 2 modes, Search and Access. The Search phase is concerned with finding the cells containing data which matches the search criteria; the Access phase is concerned with accessing those cells to extract the associated data.

For simplicity we shall assume that each cell contains a single data block of 1024 bytes stored in its memory block. If the data block is shorter than 1024 bytes, then it is simply padded out with dummy bytes. (In fact, the dummy bytes may simply be left unused, ie never read or written, in the memory block.) The data blocks also all have the same structure, which we shall take as consisting of a number of key fields Key1, Key2, Key3, etc and a number of data fields Data1, Data2, Data3, etc. Although the division of the data blocks into fields is identical for all data blocks, the number of fields, their sizes, and their locations in the data block can be chosen arbitrarily. We shall also assume that all fields are integral numbers of bytes long.

For a search, a set of key fields is defined, together with a set of search contents for those fields. Each search field is fed to the cells in turn, and for each search field, the search bytes are fed to the cells sequentially. Thus each cell receives a sequence of search bytes on the DATA bus, each accompanied on the ADD bus by its address in the data block. Further, the CAM control unit 11

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feeds a Search control signal to all cells, via the chain of logic blocks 13 in the cells.

Each time a byte address is fed to a cell in the Search phase, the corresponding location in the cell's memory block 12 is read. This reads out a byte from one of the key fields in the data block stored in the cell. The stored byte is passed to the compare unit 14, which is also fed with the search byte on the data bus DATA. (Switch 15 is disabled at this time, isolating the DATA bus from the data output of the memory block.) The compare unit 14 compares the 2 bytes fed to it, and produces a Hit or Miss signal, depending on whether the search byte and the stored byte match or don't match.

The logic block 13 of each cell contains a Match flip-flop. The CAM control unit 11 initially sets the Match flip-flops of all cells to the Hit state. In each cell, each time a comparison is made by the comparator 14 between a search byte and a stored key-field byte, the result (Hit or Miss) is fed to the Match flip-flop 16. If the stored byte and the data byte don't match, the Miss output from the comparator clears this flip-flop to Miss. After the succession of search bytes in the different key fields has been run, the Match flip-flop will remain in the Hit state only if all bytes match.

If desired, a MASK bus (not shown), with the same width as the DATA bus, can be included. Masking can time-share the DATA bus or use a second byte on the system data bus if that bus is wide enough. The MASK bus will be coupled to the comparators in all cells 10, and each bit on the MASK bus will determine whether or not the comparator compares the corresponding bits of the search and stored bytes. Thus bit-level rather than merely byte-level searching and matching can easily be implemented if desired. This allows tighter packing of the key fields if many of them are less than 1 byte long.

At the end of the Search phase, therefore, the Match flip-flop of each cell of the CAM will still be in the Hit state if a full match has been achieved for that cell, but will have been set to Miss if any failure of the matching has occurred for that cell.

Once the Search phase is finished, the Access phase follows. For this, the CAM control unit 11 sends an Access signal to the chain. At the start of this phase, there will be some unknown number of cells with their Match flip-flops

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still set at Hit. These are coupled in a logic chain through the logic blocks 13. The topmost of these cells is enabled by the Access signal and its Match flip-flop; this enabled cell disables all the following cells in the chain.

This enabled cell can be accessed over the address, data, and control buses ADD, DATA, and CONT. The address bus is coupled to the memory block 12 of the cell, and the switch 15 is enabled by the logic block 13, coupling the data bus to the memory block data path. The control bus carries a 1-bit read/write signal R/W which is, in this phase, passed through the logic block 13 to the memory block 12, to control whether reading or writing occurs. An arbitrary number of data fields can therefore be read or written, by sending the addresses of those fields to the cells in turn, and for each data field, sending the data bytes to the cells or reading them from the cells sequentially.

Once the accessing of the enabled cell is complete, the CAM control unit 11 sends a signal down the chain of cells to clear the topmost Match flip-flop to Miss. The next Match flip-flop at Hit will then enable its cell, and the data block in that cell can then be accessed. The process continues in the same way with each cell in the chain with its Match flip-flop set being enabled in turn, until all matching data blocks have been accessed.

Obviously not all data blocks found in the Search phase need be accessed in the same way. Once accessing of a block has started, the accessing may if desired be made dependant on data read from the block.

The number of matches (if any) is in general not predictable. A return line may be provided from the end of the chain of cells back to the CAM control unit 11, and arranged to change state when all Match flip-flops in the chain have been cleared to Miss.

Alternatively, the end of the sequence can be determined by software. When the end of the sequence of matching blocks is reached, any attempt to carry on reading data blocks will result in no read occurring. The data bus will therefore return its inactive state (hex-FF if it has pull-up resistors). So to determine when all matching blocks have been read, the system monitors for this data byte. To prevent genuine data bytes with the value hex-FF from being misinterpreted as the end of the set of matching blocks, a standard byte address

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is chosen in all data blocks and a non-FF-hex byte is deliberately included in that address in every data block.

Once a cell has been selected, it can of course be written as well as read. So a selected cell can simply have a complete new data block written into it, or its contents can be inspected and various parts of it changed, possibly in dependence on what it contains.

It has been assumed above that each cell contains a single data block. The optimum operating condition will be with a data block size which matches the memory block size. A data block must be confined within a single cell, so if the data blocks vary in size, the maximum size must not be larger than the cell size. However, if the data block size is considerably smaller than the cell size, it would be possible to pack 2 (or more) short data blocks into a cell. The division of the cell into such data blocks would be essentially a matter for software control. This would of course mean that a search would have to be done in 2 stages, searching first say the top halves of the cells for data block matches, and then the second halves of the cells for any further data block matches.

It is also possible for two or more data blocks to be linked together into an extended data block by suitable software techniques, so that a set of blocks so linked can be accessed as a group. A set of blocks so linked can be regarded as an extended data block whose size is not limited by the cell size. For example, each extended block can be given a unique identifier, with all cells in the block containing identical copies of the key fields and the block identifier, and each cell of the extended block containing a serial number for its position in the extended block.

Summarizing the system, therefore, it is evident that cells can be searched on any fields, and any cells matching the search can be read, modified, or completely rewritten. Although the operation was described above in terms of key and data fields, the division of the data block structure into such fields is purely notional; any bytes can be treated as key bytes, data bytes, or both.

It is not possible to access a cell by its position in the chain. However, this does not mean that a cell can become inaccessible. If a cell is to be effectively cleared, a standard byte address can be chosen for all data blocks, and filled with one data value if the data block in that cell is valid and another data

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value if the data block in the cell is cleared, ie invalid. A search for cells with the second data value in that location will then retrieve all empty data blocks in sequence. If all else fails, a search with no search fields will select all cells, which can then be written into sequentially. (That is, the Search phase is entered to set all Match flip-flops to Hit, and the Access phase is then entered immediately.)

Referring now to Fig. 3, the logic and control circuitry of the cell 10 will be described in more detail. Each cell is, as described above, coupled to the system data bus DATA and the system address bus ADD. There is also a single line chaining all the cells together. Apart from this single line, the control signals to the cells are preferably all carried on the system control bus CONT, which is coupled to all cells. In many instances, the system CPU can conveniently be arranged to generate these signals, so that this CPU constitutes the CAM control unit 11. In addition, a further line of the control bus may be coupled to the first cell as the start of the chain line through all cells. (As discussed above, a return line from the end of the chain line may usually be dispensed with.)

Fig. 3 shows the control logic unit 13 of a typical cell 10 in more detail. A memory control unit 20 is fed with three control bus signals R/W, S/A, and EN. The R/W signal is a read/write signal, which is used to determine whether the cell memory block 12 is to be read from or written to in the access phase of the operation of the CAM. The S/A signal determines whether the CAM operates in the Search or Access modes. The EN signal is an enable signal, which determines whether the CAM is enabled or not. When enabled, it is assumed that the other devices coupled to the system buses are disabled; when the CAM is disabled, other devices coupled to the system buses may be enabled without their operation being affected by the existence of the CAM.

It will be realized, of course, that the CAM (or the integrated circuit chips carrying the CAM) may be assigned a control address, so that the CAM can be enabled and disabled by means of such a control address. In effect, this means that the CAM is enabled by a suitable signal combination (the control address) on a plurality of control bus lines, but that set of control bus lines will be shared with many other devices in the system, rather than the CAM requiring a dedicated control bus line.

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The memory logic block 20 generates three output signals. There is a control signal to the bidirectional switch 15, which is used to couple the memory block 12 and the data bus DATA to the comparator 14 in the search phase and the memory block to the data bus in the access phase. There is a R/W signal to the memory block 12 to set it for reading in the search phase and for reading or writing, as the case may be, in the access phase. And there is an Enable signal to the memory block 12, to enable it in the search phase and, if the cell is the selected cell, in the access phase.

The cell also includes a chain line coupling switch 21 which couples the chain line into the cell, C-IN, to the chain line out of the cell, C-OUT. This switch in effect either couples the chain in line C-IN to the chain out line C-OUT or imposes a logic 0 on the chain out line C-OUT. It can conveniently be implemented as a pair of switches, one connected between the chain in and chain out lines and the other connecting the chain out line to earth (logic 0), with a control signal driving one switch directly and the other via an inverter. If desired, the cells may be grouped into groups with more elaborate logic for the chain line within the groups to reduce signal propagation time along the chain.

The control logic unit 13 also includes the match flip-flop 16, which has already been discussed. This flip-flop has an input from the comparator 14, as described above. It is also fed with a further signal RST from the system control bus CONT. This signal RST is a reset signal which is used to reset all match flip-flops in the CAM to the match state at the beginning of the search phase.

Finally, there is a Next control block 22. This is fed from the match flip-flop 16 and from the chain in line C-IN, and controls the chain line switch 21. If the signal on the chain in line is a logic 1 (which indicates that all cells above the present cell in the chain have been dealt with), and if the match flip-flop remains set at the end of the search phase, then the present cell is the next one to be accessed. The Next control block 22 enables the memory control block 20, and sets the switch 21 to logic 0. If the match flip-flop has been cleared, however, then the Next control block automatically disables the memory control block 20 (so that the cell cannot be accessed) and sets the switch 21 to pass the signal on the chain in line on to the chain out line.

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The Next control block is also fed with the fifth of the five signals from the system control bus, NXT, and further includes a flip-flop. If the cell is the currently selected one (ie the match flip-flop is still at match and all cells above it in the chain have been accessed), then the NXT signal effectively disables the present cell. The Next control block 21 sets the switch 21 to pass the logic 1 on to the next cell down the chain, and disables the memory control logic 20 so that the cell cannot be accessed again. The flip-flop of the Next control block is used to capture the state of the incoming C-IN signal and prevent the NXT signal from propagating down the chain. The Next control block therefore also has an output to the match flip-flop 16.

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Claims

1 A content addressable memory comprising a CAM control logic unit (11) and a plurality of cells (10) connected in a chain, each cell comprising:

- a memory block (12) coupled to a common address bus (ADD);
- a comparator (14) coupled to a common data bus (DATA) and to the data interface of the memory block (12);
- switching means (15) coupling the data interface of the memory block with the data bus, and;
- a logic block (13) including a Match flip-flop (16);

the memory being operable:

- in a Search phase to serially match a sequence of words on the common data bus (DATA) with the contents of a sequence of addresses in the memory blocks (12) of the cells (10); and
- in an Access phase, to render the cells matched in the Search phase serially available for access via the common address and data buses (ADD and DATA).

2 A content addressable memory according to claim 1 wherein each cell contains a memory block (12), a logic block (13), a comparator (14), and a bidirectional switch (15).

3 A content addressable memory according to either previous claim implemented on an integrated circuit chip.

4 A content addressable memory according to claim 3 wherein several such chips can be chained.

5 A content addressable memory according to claim 4 wherein each chip includes a control unit which can be disabled.

6 A content addressable memory according to any previous claim including a MASK bus input which determines which bits of the words of the sequence of words are used for matching in the Search phase.

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7 A content addressable memory according to any previous claim including a return line from the end of the chain of cells back to the CAM control unit 11 which changes state when all Match flip-flops in the chain have been accessed.

8 A method of operating a content addressable memory according to any previous claim wherein a standard byte address is chosen in all data blocks and a byte different from the inactive state of the data bus is included in that address in every data block.

9 A method of operating a content addressable memory according to any previous claim wherein each cell is divided into a plurality of distinct data blocks.

10 A method of operating a content addressable memory according to any previous claim wherein a plurality of cells are combined into an extended data block with all cells of the block containing corresponding key fields.

11 A method of operating a content addressable memory according to any previous claim wherein a standard byte address is chosen in all data blocks and filled with one data value if the data block in that cell is valid and another data value if the data block in the cell is cleared, ie invalid.

12 Any novel and inventive feature or combination of features specifically disclosed herein within the meaning of Article 4H of the International Convention (Paris Convention).

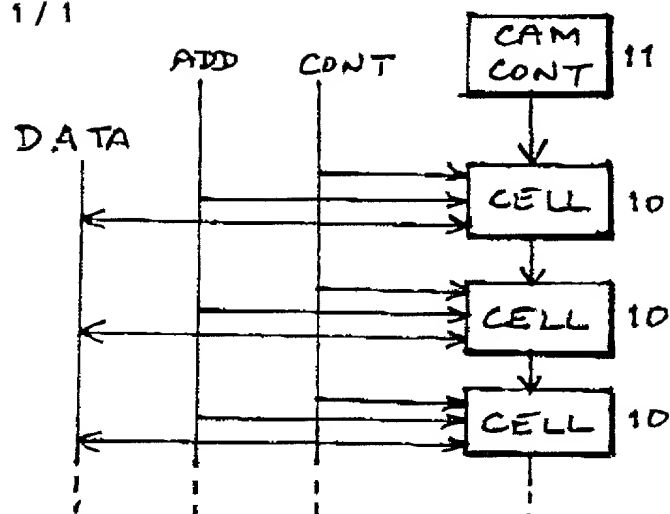
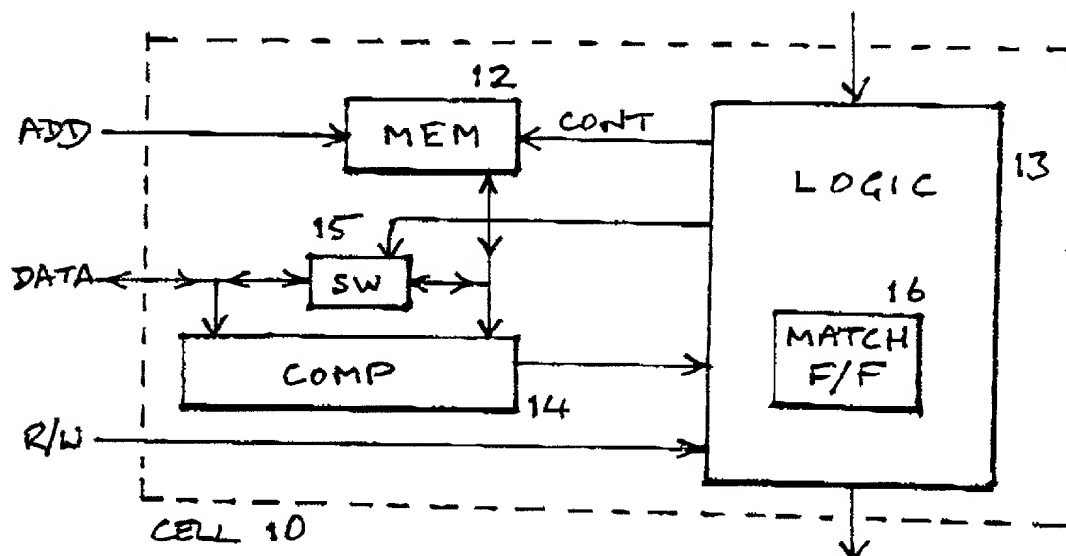
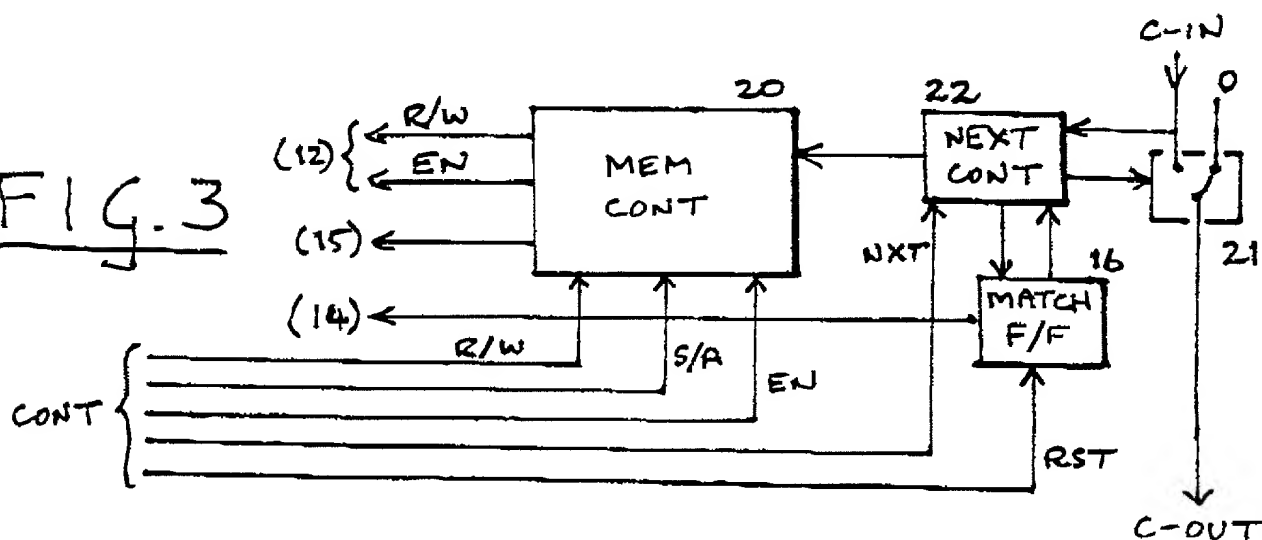
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FIG. 1FIG. 2FIG. 3

DECLARATION FOR PATENT APPLICATION AND APPOINTMENT OF ATTORNEY

As the below-named inventor, I hereby declare that my residence, post office address and citizenship are as stated below next to my name, I believe that I am the original first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention (Design, if applicable) entitled
CONTENT ADDRESSED MEMORIES

the specification of which (check one):

- ☐ is attached hereto.
- ☒ was filed on 15 January 2001 as Application Serial No 09/743,713; and
- ☒ was filed on 15 July 1999 as International Application (PCT) No PCT/GB99/02288, and was amended on _____ (if applicable)

I hereby state that I have reviewed and understand the contents of the above-identified specification, including the claims, as amended by any amendment(s) referred to above. I acknowledge the duty to disclose information which is material to the examination of this application in accordance with *Title 37, Code of Federal Regulations, § 1.56*. I hereby claim foreign priority benefits under *Title 35, United States Code, § 119* of any foreign application(s) for patent or inventor's certificate listed below and have also identified below any foreign application for patent or inventor's certificate having a filing date before that of the application on which the priority is claimed.

PRIOR FOREIGN APPLICATION(S)

NUMBER	COUNTRY	DAY/MONTH/YEAR FILED	PRIORITY CLAIMED
9815370.3	Great Britain	16 July 1998	<input checked="" type="checkbox"/> Yes <input type="checkbox"/> No
			<input type="checkbox"/> Yes <input type="checkbox"/> No

I hereby claim the benefit under *Title 35, United States Code, § 120* of any United States application(s) or PCT international application(s) designating The United States of America listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in that/those prior application(s) in the manner provided by the first paragraph of *Title 35, United States Code, § 112*, I acknowledge the duty to disclose material information as defined in *Title 37, Code of Federal Regulations, § 1.56* which occurred between the filing date of the prior application(s) and the national or PCT international filing date of this application:

APPLICATION NUMBER	FILING DATE	STATUS (Patented, Pending or Abandoned)

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine, or imprisonment, or both, under *Section 1001 of Title 18 of the United States Code*, and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

POWER OF ATTORNEY: I (We) hereby appoint as my (our) attorneys, with full powers of substitution and revocation, to prosecute this application and transact all business in the Patent and Trademark Office connected therewith: Allan M. Lowe, Registration Number 19,641; Benjamin J. Hauptman, Registration Number 29,310; Michael G. Giltman, Registration Number 19,114; Kenneth M. Berger, Registration Number 37,093; and Randy A. Noranbrock, Registration Number 42,940.

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I hereby authorize the U.S. attorneys and agents named herein to accept and following instructions from MG Harman as to any actions to be taken in the U.S. Patent and Trademark Office regarding this application without direct communication between the U.S. attorneys and the undersigned. In the event of a change in the person(s) from whom instructions may be taken, the U.S. attorneys will be so notified by the undersigned.

☐ See following page(s) for additional joint inventors.

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